

IN THE SPECIFICATION

Please amend the title as follows:

METHOD AND [[AN]] ARRANGEMENT FOR MANAGING PACKET QUEUES IN SWITCHES

Please amend the paragraph at page 1, line 18 as follows:

State of the art Background

It is previously known to divide data streams for various reasons. In the Japanese published document number JP 59-103147, an A/D converter is shown having two parallel buffers. Data given from the A/D converter is divided to be stored alternately in one of the buffers depending on an occupancy of the buffer. The Japanese published document number JP 11-008631 shows an ATM cell transmission flow control system having a divided buffer. The Japanese published document number JP 03-100783 shows a queue buffer system including a queue buffer and an external memory. When the queue buffer is filled up with tokens, tokens overflowing the queue buffer are written in the external memory.

Please amend the paragraph at page 3, line 22 as follows:

With reference to figure 2, ~~we now look at the data flows belonging to one output port are~~ shown. To the left there is an incoming data stream (A+B) which is larger than the output capacity D (e.g. 1 Gbit/s) of the output port. A basic concept of the invention is to divert only a part of the incoming data stream when it becomes larger than the output capacity of the output port. Thus, it may be seen that a first part A of the data stream is sent to the internal queue of the primary memory and a second part B is sent to the external memory (via the store queue 6). The first part A may be selected a little smaller than the output capacity D of the output port, so that a small data flow C may be integrated back from the external memory to the internal queue for reasons that will be explained below. The division of the data stream results in that the capacity of the output port is always ~~utilised~~ utilized to the largest possible extent. The output port will not be blocked by diverting the whole data stream to the external memory.

Please amend the paragraph at page 4, line 19 as follows:

Each priority group may also be divided into subgroups to achieve even finer granularity. The finer granularity, the more closely the part A to be sent directly to the internal queue may be adapted. In this example each priority group is divided into four so-called hash groups. The hash groups are formed by means of ~~other~~ criteria other than priority. In a preferred embodiment of the invention, a hash group is formed by looking at a part of an arriving data packet and calculating a value based on that part, so that the packets will be evenly distributed in four groups, provided that the data parts are randomly distributed. Suitably, flow information is used that is constant during a session, e.g. an originating or destination address part of the data packet. This will result in that there is a logical continuity within the hash groups.